#### Saturation

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## Formal methods

- Safety critical and embedded systems

   Railway, automotive industry, air transportation
   Reliability is an important issue
- Design time analysis
  - These models can be used for implementation



## Model checking

#### Automatic verification method



- Prerequisite:
  - Exploring and representing the reachable states
- Problem:
  - State space explosion
  - Time and space requirements



## Saturation algorithm



- Efficient solution for:
  - State space generation
  - Model checking
- Symbolic algorithm
  - Encoding of states
  - Special underlying data structures
    - Multi Valued Decision Diagrams (MDD-s)
- Special iteration strategy
  - Efficient for asynchronous models



#### Multi Valued Decision Diagrams

- Derived from decision trees

   variables are ordered into levels
- Example:
  - o only binary variables





#### Multi Valued Decision Diagrams

- Derived from decision trees
  - variables are ordered into levels
- Special reduction rules
  - in a bottom-up fashion, applying reduction from level-to-levels
- Compact representation of multi valued functions



## Symbolic algorithm

- Symbolic encoding instead of explicit state representation
  - Decomposition is needed
- Saturation uses component wise encoding





## Special iteration

- Local exploration in a greedy manner
- Exploring global synchronization events if needed
- Uses the primarily defined order of the decision diagram variable encoding
- Efficient for Globally Asynchronous, Locally Synchronous models (GALS)





#### State space representation with MDDs



MÚEGYETE



#### State space representation with MDDs



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- Dining philosophers
   5 philosophers
- State space representation
  - o 1364 states
  - o 19 nodes







- Dining philosophers
   0 10 philosophers
- State space representation
  - o 1,860,498 states
  - o 40 nodes







- Dining philosophers
   20 philosophers
- State space representation
  - 3,461,452,808,002
     states
  - o 80 nodes





- Slotted Ring communication protocol

   2 slots
- State space representation
  - 52 states
  - o 14 nodes







- Slotted Ring communication protocol

   4 slots
- State space representation
  - o 5136 states
  - o 30 nodes







- Slotted Ring communication protocol

   8 slots
- State space representation
  - o 68,026,624 states
  - 103 nodes



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- Slotted Ring communication protocol
   20 slots
- State space representation
  - $\circ$  10<sup>20</sup> states

 487 nod
 Scales up to about 200 slots in the ring and about 10<sup>200</sup> states
 (60902 nodes in the state space MDD, the full state space generation lasted 222 seconds long)





- Flexible manufacturing system
   5 item
- State space representation o about 2,900,000 states o 248 nodes Not so nice, but still efficient 🙂



- Tower of Hanoi game
  - It consists of three rods, and a number of disks of different sizes which can slide onto any rod.
  - o Rules:
    - Only one disk may be moved at a time
    - Each move consists of taking the upper disk from one of the rods and sliding it onto another rod, on top of the other
    - No disk may be placed on top of a smaller disk

Synchronous model:

 at most 4 transitions are enabled from each state







- Tower of Hanoi game
   0 12 disks
- State space representation
   531 441 (3<sup>12</sup>) states
  - o 12 nodes

#### Unfortunately:

- during the exploration we construct more nodes
- the state space generation took 58 seconds
- huge number of transitions in the model



- Tower of Hanoi game
   0 12 disks
- State space representation
  - 531 441 (3<sup>12</sup>) states
  - o 12 nodes

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#### Unfortunately:

#### **Conclusion:**

**Efficient state space representation** 

#### **Efficient iteration**

- For asynchronous models -



125 level: 12

## Problems

- Efficiency of the algorithm highly depends:
   Decomposition
  - Variable ordering
    - Bottleneck of symbolic methods

 Best performance if this information is provided manually





#### Motivation for bounded model checking







#### Motivation for bounded model checking



#### **Bounded model checking (BMC)**

- explores a k-bounded part of the state space (usually in a breadth first manner)
- examines the specification on this smaller part





#### **Bounded Saturation**

- Bounded model checking
  - explores a k-bounded part of the state space
    - usually in a breadth first manner
  - examines the specification on this smaller part
- Saturation
  - Explores the state space in an irregular recursive order
  - Difficult to bound the exploration
  - There is no distance information in the MDD-s





#### **Bounded Saturation**

- Bounded model checking
  - explores a k-bounded part of the state space
    - usually in a breadth first manner
  - examines the specification on this smaller part
- Saturation

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- Explores the state space in an irregular recursive order
  - New data structure:
  - Edge Valued Decision Diagrams (EDDs)
    - MDD based data structure enriched with distance information





## Saturation Based bounded model checking



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#### Automata theoretic model checking

Verifying  $\omega$ -regular properties





#### Traditional approach





- Can be obtained by:
  - Complementing the Büchi-automaton describing the specification (not always possible)
  - Compiling the negation of the specification described in linear temporal logic
- Compilation algorithm and optimization is crucial
  - It is the *"*exponential part" of the model checking process
  - Product computation complexity is linear in the size of the automaton





- Can be obtained from higher level models by the means of statespace generation
  - State space can be huge, especially in asynchronous systems
- Symbolic representation can solve storage problems
  - But it complicates the generation process
- Saturation algorithm provides efficient symbolic state space exploration
  - Especially in asynchronous systems...





- Product state space is the Cartesian product of the sets of states of the automata
  - But only a part is reachable from the initial state
- A word of the product language is an accepting trace of the product automaton
  - An accepting trace is a reachable loop that contains accepting states
- Task: find such a loop or prove that there does not exist one
  - o Efficient on-the-fly algorithms on explicit graph representations
  - Symbolic approaches are cumbersome







- Compilation algorithm handling linear temporal logics (LTL)
  - Past and future time linear temporal expressions
- Heavy optimization and minimization of the specification automaton
  - Reduction of the temporal expressions
  - Removing redundancy from the automaton

# A compiled and optimized automaton can be used multiple times!





- Product state space is generated together with the system state space
  - Using a modified version of the saturation algorithm
- Emptiness checking is performed simultaneously during the generation process
  - Unlike traditional symbolic approaches, accepting loops are being detected incrementally
  - State space generation halts immediately when a loop is found
  - This gives us an incremental, on-the-fly LTL model checking process
    - $\circ$  Supporting LTL, PLTL and any  $\omega$ -regular specification properties



